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ABSTRACT

A two-coloring of the vertices V of the hypergraph $H = (V, E)$ by red and blue has discrepancy d if d is the largest difference between the number of red and blue points in any edge. Let $f(n)$ be the fewest number of edges in an n -uniform hypergraph without a coloring with discrepancy 0. Erdős and Sós asked: is $f(n)$ unbounded?

N. Alon, D. J. Kleitman, C. Pomerance, M. Saks and P. Seymour [1] proved upper and lower bounds in terms of the smallest non-divisor (snd) of n (see (1)). We refine the upper bound as follows:

$$f(n) \leq c \log \text{snd } n.$$

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1. Introduction

A hypergraph is a pair (V, E) , where V is a finite set whose elements are called vertices and E is a family of subsets of V , called edges. A hypergraph is n -uniform if every edge has size n . A vertex 2-coloring of a hypergraph (V, E) is a map $\pi : V \rightarrow \{1, 2\}$.

The *discrepancy* of a coloring is the maximum over all edges of the difference between the number of vertices of two colors in the edge. The *discrepancy* of a hypergraph is the minimum discrepancy of a coloring of this hypergraph. The general discrepancy theory is set out in [2,6,4].

Let $f(n)$ be the minimal number of edges in an n -uniform hypergraph (all edges have size n) having positive discrepancy. Obviously, if $2 \nmid n$ then $f(n) = 1$; if $2 \mid n$ but $4 \nmid n$ then $f(n) = 3$. Erdős and Sós asked whether $f(n)$ is bounded or not. N. Alon, D. J. Kleitman, C. Pomerance, M. Saks and P. Seymour [1] proved the following Theorem, showing in particular that $f(n)$ is unbounded.

Theorem 1.1. *Let n be an integer such that $4 \mid n$. Then*

$$c_1 \frac{\log \text{snd}(n/2)}{\log \log \text{snd}(n/2)} \leq f(n) \leq c_2 \frac{\log^3 \text{snd}(n/2)}{\log \log \text{snd}(n/2)}, \quad (1)$$

where $\text{snd}(x)$ stands for the least positive integer that does not divide x .

To prove the upper bound they introduced several quantities. Let \mathcal{M} denote the set of all matrices M with entries in $\{0, 1\}$ such that the equation $Mx = e$ has exactly one non-negative solution (here e stands for the vector with all entries equal to 1). This unique solution is denoted x^M . Let $z(M)$ be the least integer such that $z(M)x^M$ is integer and let $y^M = z(M)x^M$. For each positive integer n , let $t(n)$ be the least r such that there exists a matrix $M \in \mathcal{M}$ with r rows such that $z(M) = n$ (obviously, $t(n) \leq n+1$ because $z(J_{n+1} - I_{n+1}) = n$, where J_{n+1} is the $(n+1) \times (n+1)$ matrix with unit entries; I_{n+1} is the $(n+1) \times (n+1)$ identity matrix). The upper bound in (1) follows from the inequality $f(n) \leq t(m)$ for such m that $\lfloor \frac{n}{m} \rfloor$ is odd.

Then N. Alon and V. H. Vū [3] showed that $t(m) \leq (2 + o(1)) \frac{\log m}{\log \log m}$ for infinitely many m . However they marked that trueness of inequality $t(m) \leq c \log m$ for arbitrary m is not clear.

Our main result is the following

Theorem 1.2. *Let n be a positive integer number. Then*

$$f(n) \leq c \log \text{snd}(n), \quad (2)$$

for some constant $c > 0$.

Corollary 1.3. *Let n be a positive integer number. Then*

$$f(n) \leq c \log \log n,$$

for some constant $c > 0$.

The construction of the hypergraph with positive discrepancy which yields Theorem 1.2 uses a matrix with determinant $\text{snd}(n)$ and small entries satisfying some additional technical properties. Before coming to a general construction we give an example with a specific 2×2 matrix which shows the vague idea.

2. Example

Example 2.1. Let us consider the matrix $A = \begin{pmatrix} 3 & 5 \\ 1 & 8 \end{pmatrix}$ and suppose that n is not divisible on $\det A = 19$. Consider the system

$$A \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} n \\ n+t \end{pmatrix}. \quad (3)$$

The solution of the system is $a = (3n - 5t)/19$, $b = (2n + 3t)/19$, which is integral if and only if $t = 12n \pmod{19}$ i.e. t has prescribed residue modulo 19. Since n is not divisible on 19, t is not equal to zero modulo 19. So one can choose $-19 < t < 19$ such that t has prescribed residue modulo 19 and t is odd. Also, assume that $n/8 > t > -2n/3$ which is certainly true if $n > 200$. Then a and b are positive and also $b > t$ and a, b tend to infinity simultaneously with n .

Let us construct an n -uniform hypergraph H with positive discrepancy. Consider disjoint vertex sets A_1, A_2, A_3 of size a and B_1, \dots, B_8 of size b . If $t < 0$ then consider a vertex set T of size $|t|$ and set $C := B_1 \cup T$; if $t > 0$ let T be a t -vertex subset of B_1 and define $C := B_1 \setminus T$. The edges of H are listed:

$$\begin{aligned} & A_1 \cup A_2 \cup A_3 \cup B_1 \cup B_2 \cup B_3 \cup B_4 \cup B_5 \\ & A_1 \cup A_2 \cup A_3 \cup B_1 \cup B_2 \cup B_3 \cup B_4 \cup B_6 \\ & A_1 \cup A_2 \cup A_3 \cup B_1 \cup B_2 \cup B_3 \cup B_4 \cup B_7 \\ & A_1 \cup A_2 \cup A_3 \cup B_1 \cup B_2 \cup B_3 \cup B_4 \cup B_8 \\ & A_1 \cup A_2 \cup A_3 \cup B_2 \cup B_3 \cup B_4 \cup B_5 \cup B_8 \\ & A_1 \cup A_2 \cup A_3 \cup B_1 \cup B_3 \cup B_4 \cup B_5 \cup B_8 \\ & A_1 \cup A_2 \cup A_3 \cup B_1 \cup B_2 \cup B_4 \cup B_5 \cup B_8 \\ & A_1 \cup A_2 \cup A_3 \cup B_1 \cup B_2 \cup B_3 \cup B_5 \cup B_8 \\ & A_1 \cup C \cup B_2 \cup B_3 \cup B_4 \cup B_5 \cup B_6 \cup B_7 \cup B_8 \end{aligned}$$

$$A_2 \cup C \cup B_2 \cup B_3 \cup B_4 \cup B_5 \cup B_6 \cup B_7 \cup B_8$$

$$A_3 \cup C \cup B_2 \cup B_3 \cup B_4 \cup B_5 \cup B_6 \cup B_7 \cup B_8.$$

Obviously, if H has a coloring with discrepancy 0, then $d(B_5) = d(B_6)$, where $d(X)$ is the difference between blue and red vertices in X , because the second edge can be reached by replacing B_5 on B_6 in the first edge. Similarly one can deduce that $d(A_i) = d(A_j)$ and $d(B_i) = d(B_j)$ for all pairs i, j . So one can put $k := d(A_i), l := d(B_i)$. Because of the first edge we have $3k + 5l = 0$. Obviously, k and l are odd numbers, so the minimal solution is $k = 5, l = -3$ (or $k = -5, l = 3$ which is the same because of red-blue symmetry). But then the last edge gives $|k + 8l| \leq |t|$ which contradicts with $|k + 8l| \geq 19 > |t|$.

So we got an example if $19 \nmid n$ and $n > 200$ of an n -uniform hypergraph with 11 edges and positive discrepancy.

The number of edges in this example equals $11 = 3 + 8$, the sum of maximal entries in the columns of A . This is essentially (up to multiplicative constant) the general property of our construction.

3. Proofs

Proof of Theorem 1.2. Let us denote $\text{snd}(n)$ by q . We should construct a hypergraph with at most $c \log q$ edges and positive discrepancy. Take m such that $2^m - 1 \leq q \leq 2^{m+1} - 2$. Then

$$q - (2^m - 1) = \sum_{i=0}^{m-1} \varepsilon_i 2^i \quad \text{for some } \varepsilon_i \in \{0, 1\},$$

therefore

$$q = \sum_{i=0}^{m-1} \eta_i 2^i, \quad \text{where } \eta_i = 1 + \varepsilon_i \in \{1, 2\}.$$

Consider m vectors in \mathbb{Z}^m :

$$v_0 = (\eta_0, \dots, \eta_{m-1}),$$

$$v_i = (\eta_0, \dots, \eta_{i-2}, \eta_{i-1} + 2, \eta_i - 1, \eta_{i+1}, \dots, \eta_{m-1}) \quad \text{for } i = 1, \dots, m - 1, \text{ i.e.}$$

$$v_{i,k} = \begin{cases} \eta_k, & k \neq i, i - 1 \\ \eta_k - 1, & k = i \\ \eta_k + 2, & k = i - 1. \end{cases}$$

Note that the vector $u = (1, 2, \dots, 2^{m-1})$ satisfies a system of linear equations

$$\langle v_i, u \rangle = q; \quad i = 0, \dots, m - 1.$$

Assume that q is odd. Choose odd $\delta \in (-q, q)$ such that $x_0 := \frac{n+\eta_{m-1}\delta}{q}$ is integer. Define

$$x_i := 2^i x_0 \text{ for } i = 1, \dots, m-2; \quad x_{m-1} := 2^{m-1} x_0 - \delta,$$

then the vector $x = (x_0, \dots, x_{m-1})$ satisfies $\langle v_i, x \rangle = n$ for $i = 0, \dots, m-2$, $\langle v_{m-1}, x \rangle = n + \delta$.

In the case $q = 2^m \geq 8$ we have $n \equiv 2^{m-1} \pmod{q}$ and $\eta_0 = 2, \eta_1 = \dots = \eta_{m-1} = 1$.

Choose $x = (x_0, \dots, x_{m-1})$ so that $\langle v_1, x \rangle = \langle v_{m-1}, x \rangle = n + 1$ and $\langle v_i, x \rangle = n$ for $i = 0, 2, 3, \dots, m-2$. The solution is given by

$$x_0 := \frac{n + 2^{m-1}}{q}; \quad x_1 := 2x_0 - 1; \quad x_i := 2^{i-1}x_1 \text{ for } i = 2, \dots, m-2; \quad x_{m-1} := 2^{m-2}x_1 - 1.$$

Now let us construct a hypergraph in the following way: for $i = 0, \dots, m-1$ let us take 4 sets A_i^j ($j = 1, \dots, 4$) of vertices of size x_i such that all $4m$ sets A_i^j are disjoint. Let the edge e_0 be the union of A_i^j over $0 \leq i \leq m-1$ and $1 \leq j \leq \eta_i$. By the choice of x_i and η_i we have $|e_0| = n$. Then we add an edge

$$\bigcup_{0 \leq i \leq m-1} \bigcup_{\substack{1 \leq j \leq \eta_i \text{ for } i \neq k \\ j \in R \text{ for } i = k}} A_i^j$$

for every k and for every $R \subset [4]$ such that $|R| = \eta_k$. Clearly there are at most $6m$ such edges. Let us say that they form *the first collection of edges*. Finally, for every $1 \leq k \leq m-1$ we add the edge

$$\bigcup_{0 \leq i \leq m-1} \bigcup_{\substack{1 \leq j \leq \eta_i \text{ for } i \neq k, k-1 \\ 1 \leq j \leq \eta_i + 2 \text{ for } i = k-1 \\ 1 \leq j \leq \eta_i - 1 \text{ for } i = k}} A_i^j,$$

which form *the second collection of edges*.

Summing up we have hypergraph with at most $7m$ edges; at most 2 of them have size not equal to n . Let us correct these edges in the simplest way: if an edge has size less than n then we add arbitrary vertices; if an edge has size greater than n then we exclude arbitrary vertices.

Suppose that our hypergraph has discrepancy 0, so it has a proper coloring π . For every set A_i^j denote by $d(A_i^j)$ the difference between the numbers of red and blue vertices of π in A_i^j . Obviously, $d(A_i^{j_1}) = d(A_i^{j_2})$ because there are edges e_1, e_2 from the first collection such that e_2 can be obtained from e_1 by the replacement of $A_i^{j_1}$ to $A_i^{j_2}$. So we may write d_i instead of $d(A_i^j)$.

If q is odd then the vector $d = (d_0, \dots, d_{m-1})$ satisfies

$$\langle v_i, d \rangle = 0 \text{ for } i = 0, 1, \dots, m-2 \text{ and } \langle v_{m-1}, d \rangle = s$$

for some odd $s \in (-q, q)$. Considering consequent differences of these equations we get

$$d_i = 2^i d_0 \text{ for } i = 0, \dots, m-2; \quad d_{m-1} = 2^{m-1} d_0 - s; \quad 0 = \sum \eta_i d_i = d_0 q - \eta_{m-1} s,$$

which fails modulo q . A contradiction. In the case $q = 2^m$ we get a similar contradiction, as $(2^{m-1} - 1) \pm 1$ is not divisible by 2^m .

Thus we get a hypergraph on at most $7m = O(\log q)$ edges with positive discrepancy, the claim is proven. \square

4. Discussion

- In fact, during the proof we have constructed a matrix of size of $O(\log k)$ with bounded integer coefficients and with determinant $k := \text{snd}(n)$. By Hadamard inequality, the determinant k of $m \times m$ matrix with bounded coefficients satisfies $k = O(\sqrt{m})^m$, thus $\log k = O(m \log m)$, $m \geq \text{const} \cdot \log k / \log \log k$. We suppose that actually a matrix of size $O(\log k / \log \log k)$ with bounded integer coefficients and determinant k always exists; and moreover, it may be chosen satisfying additional properties which allow to replace the main estimate with $f(n) \leq c \log \text{snd}(n) / \log \log \text{snd}(n)$ (which asymptotically coincides with the lower bound).
- It turns out, that for a fixed value of $q = \text{snd}(n)$ and some values of n modulo q , a hypergraph, constructions of above type have the discrepancy separated from zero. In particular, in Example 2.1 the choice $n \in \{\pm 4, \pm 7\}$ modulo 19 leads to the discrepancy 6.
- For fixed r and large enough n using Theorem 1.2 one can construct an n -uniform hypergraph with discrepancy at least r and $O(\ln \text{snd} [n/r])^r$ edges (here $[x]$ stands for the nearest integer to x), as follows: let H_0 be a hypergraph realizing $f([n/r])$, H_1, \dots, H_{2r-1} be vertex-disjoint copies of H . Let $V := V(H_1) \sqcup \dots \sqcup V(H_{2r-1})$, $E := \{\sqcup e_i \mid i \in A \subset [2r-1], |A| = r\}$. By the construction, every H_i has discrepancy at least 2; so by pigeonhole principle (V, E) has discrepancy at least $2r$. Define $l := r[n/r] - n$. Finally, if $l > 0$, then exclude arbitrary l vertices from every edge $e \in E$; else add arbitrary l vertices to every edge $e \in E$; denote the result by H . By definition $l \leq r$, so the discrepancy of H is at least r . Since $|E(H_i)| = f([n/r])$, we have

$$|E(H)| = \binom{2r-1}{r} f([n/r])^r = O(\ln \text{snd} [n/r])^r \leq O(\ln \ln n)^r.$$

- A. Raigorodskii independently asked the same question in a more general form: he introduced the quantity $m_k(n)$ that is the minimal number of edges in a hypergraph without a vertex 2-coloring such that every edge has at least k blue vertices and at least k red vertices. So $m_k(n)$ is the minimal number of edges in a hypergraph with discrepancy at least $n - 2k + 2$, in particular $f(n) = m_{n/2}(n)$ for even n .

For the history and the best known bounds on $m_k(n)$ see [7]. Note that our result replaces the bound $m_k(2k+r) = O(\ln k)^{r+1}$ [5] with $m_k(2k+r) = O(\ln \ln k)^{r+1}$ for a constant r . It worth noting, that in the case $n = O(k)$ the behavior of $m_k(n)$ is completely unclear.

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